Department of Computer Science



### University Institute of Engineering DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Bachelor of Engineering Subject Name: System Programming Subject Code: CST-315

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Compilers





# Chapter-1.2 Assembler

- Multi-Pass Assemblers
- Advanced Assembly Process

# Forward Reference

- □ All symbol-defining directives do *not* allow forward reference for 2-pass assembler
  - e.g., EQU, ORG...
  - All symbols used on the *right-hand side* of the statement must have been defined previously
  - E.g. (Cannot be assembled in 2-pass assm.) BETA EQU DELTA DELTA RESW 1

## 2.4 Assembler Design Options

One-pass assemblers

Multi-pass assemblers

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## 2.4.1 One-Pass Assemblers

- □ Goal: avoid a second pass over the source program
- □ Main problem
  - Forward references to *data items* or *labels on instructions*
- □ Solution
  - Data items: require all such areas be defined before they are referenced
  - Label on instructions: cannot be eliminated
    - E.g. the logic of the program often requires a forward jump
    - □ It is too inconvenient if forward jumps are not permitted

### Two Types of One-Pass Assemblers:

- □ <u>Load-and-go</u> assembler
  - Produces object code directly in memory for immediate execution

- □ The other assembler
  - Produces usual kind of object code for later execution

# Load-and-Go Assembler

- No object program is written out, no loader is needed
- □ Useful for program development and testing
  - Avoids the overhead of writing the object program out and reading it back in
- Both one-pass and two-pass assemblers can be designed as load-and-go
  - However, one-pass also avoids the overhead of an additional pass over the source program
- □ For a load-and-go assembler, the actual address must be known at assembly time.

### Forward Reference Handling in One-pass Assembler

- When the assembler encounter an instruction operand that has not yet been defined:
  - 1. The assembler omits the translation of operand address
  - 2. Insert the symbol into SYMTAB, if not yet exist, and mark this symbol *undefined*
  - 3. The address that refers to the undefined symbol is added to <u>a list of</u> <u>forward references</u> associated with the symbol table entry
  - 4. When the definition for a symbol is encountered
    - 1. The forward reference list for that symbol is scanned
    - 2. The proper address for the symbol is inserted into any instructions previous generated.

### Handling Forward Reference in One-pass Assembler (Cont.)

- □ At the end of the program
  - Any SYMTAB entries that are still marked with \* indicate <u>undefined symbols</u>
    - □ Be flagged by the assembler as errors
  - Search SYMTAB for the symbol named in the END statement and jump to this location to begin execution of the assembled program.

# Sample Program for a One-Pass Assembler (Fig. 2.18)

Line	Loc	Source statement			Object code		
0 1 2 3 4 5 6 9	1000 1000 1003 1006 1009 100C 100F	COPY EOF THREE ZERO RETADR LENGTH BUFFER	START BYTE WORD WORD RESW RESW RESB	1000 C'EOF' 3 0 1 1 4096	454F46 000003 000000		
9 10 15 20 25 30 35 40 45 55 60 55 60 50 70 75 110	200F 2012 2015 2018 201B 2021 2021 2024 2027 202A 202D 2030 2033 2036	CLOOP ENDFIL	STL JSUB LDA COMP JEQ JSUB J LDA STA LDA STA JSUB LDL RSUB	RETADR RDREC LENGTH ZERO ENDFIL WRREC CLOOP EOF BUFFER THREE LENGTH WRREC RETADR	$\begin{array}{c} 141009\\ 48203D\\ 00100C\\ 281006\\ 302024\\ 482062\\ 302012\\ 001000\\ 0C100F\\ 001003\\ 0C100C\\ 482062\\ 081009\\ 4C0000\\ \end{array}$		

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# Sample Program for a One-Pass Assembler (Fig. 2.18) (Cont.)

110		۲	<u>CIEDO</u>	TINE TO READ	RECORD INTO BUFFI
115			SUBROU	TINE TO READ	RECORD INTO BUFFI
120		•			2
121	2039	INPUT	BYTE	X'F1'	<b>F</b> 1
122	203A	MAXLEN	WORD	4096	001000
124					
125	203D	RDREC	LDX	ZERO	041006
130	2040		LDA	ZERO	001006
135	2043	RLOOP	TD	INPUT	E02039
140	2046		JEQ	RLOOP	302043
145	2049		RD	INPUT	D82039
150	204C		COMP	ZERO	281006
155	204F		JEQ	EXIT	30205B
160	2052		STCH	BUFFER,X	54900F
165	2055		TIX	MAXLEN	2C203A
170	2058		$_{\rm JLT}$	RLOOP	382043
175	205B	EXIT	STX	LENGTH	10100C
180	205E		RSUB		4C0000
195					

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# Sample Program for a One-Pass Assembler (Fig. 2.18) (Cont.)

тар		*			
200		•	SUBROU	TINE TO WRITE	RECORD FROM BUFFER
205					
206	2061	OUTPUT	BYTE	X'05'	05
207		•			
210	2062	WRREC	LDX	ZERO	041006
215	2065	WLOOP	TD	OUTPUT	E02061
220	2068		JEQ	WLOOP	302065
225	206B		LDCH	BUFFER,X	50900F
230	206E		WD	OUTPUT	DC2061
235	2071		TIX	LENGTH	2C100C
240	2074		JLT	WLOOP	382065
245	2077		RSUB		4C0000
255			END	FIRST	

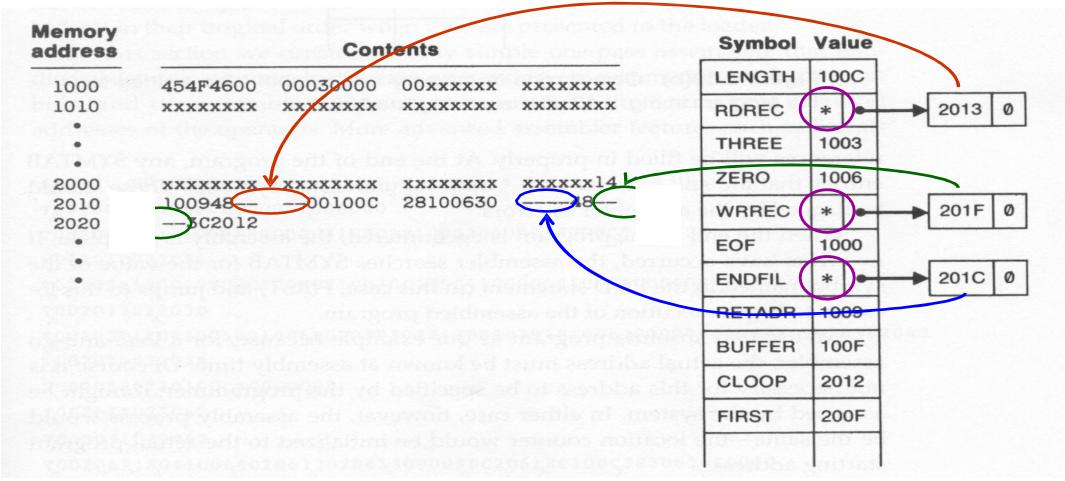
Figure 2.18 Sample program for a one-pass assembler.

# Example

- **•** Fig. 2.19 (a)
  - Show the object code in memory and symbol table entries after scanning line 40
  - Line 15: forward reference (RDREC)
    - □ Object code is marked ----
    - □ Value in symbol table is marked as \* (undefined)
    - □ Insert *the address of operand* (2013) in a list associated with RDREC
  - Line 30 and Line 35: follow the same procedure

### Object Code in Memory and SYMTAB

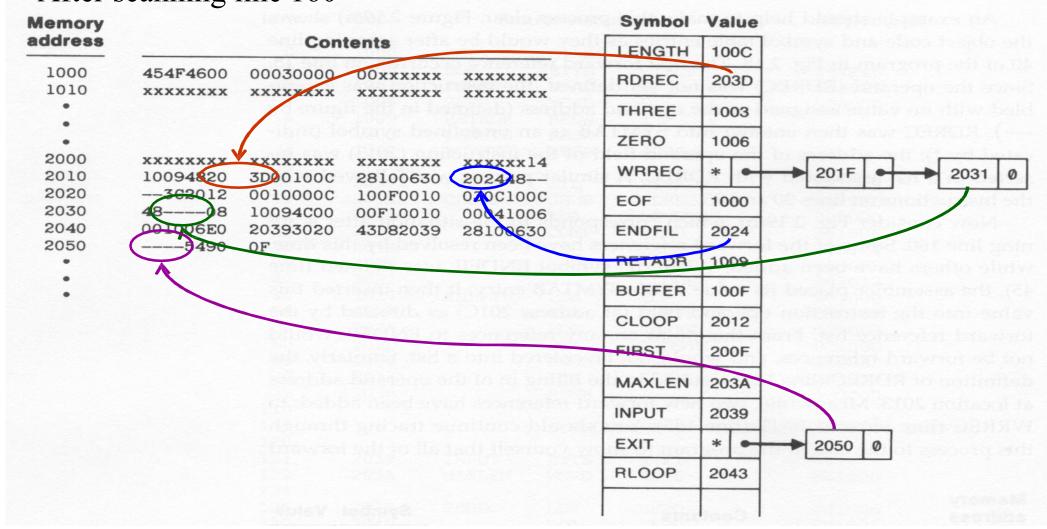
#### After scanning line 40



# Example (Cont.)

- □ Fig. 2.19 (b)
  - Show the object code in memory and symbol table entries after scanning line 160
  - Line 45: ENDFIL was defined
    - □ Assembler place its value in the SYMTAB entry
    - Insert this value into the address (at 201C) as directed by the forward reference list
  - Line 125: RDREC was defined
    - **Follow the same procedure**
  - Line 65 and 155
    - □ Two new forward reference (WRREC and EXIT)

#### Object Code in Memory and SYMTAB After scanning line 160



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### Object Code in Memory and SYMTAB Entries for Fig 2.18 (Fig. 2.19b)

Memory					Symbol	Valu	ш
address		Cont	tents		LENGTH	100	C
1000	454F4600	00030000	00xxxxxx	******	RDREC	203	Ð
1010	XXXXXXXX	XXXXXXXXX	XXXXXXXX	*****	THREE	100	3
-					ZERO	100	6
2000	*******	XXXXXXXX 3D00100C	XXXXXXXX 28100630	xxxxxx14 202448	WRREC	*	
2010 2020	10094820 302012	0010000C	100F0010	030C100C	EOF	100	)0
2030 2040	4808 001006E0	10094C00 20393020	00F10010 43D82039	00041006 28100630	ENDFIL	202	24
2050	5490	OF			RETADR	100	29
•					BUFFER	100	DF
•					CLOOP	20	12
					FIRST	200	O F
					MAXLEN	203	3/
					INPUT	20:	39
					EXIT	*	•
					RLOOP	20	4:
						-	

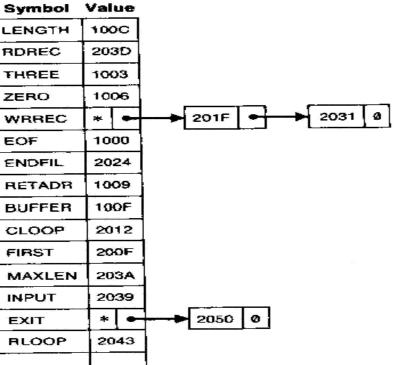


Figure 2.19(b) Object code in memory and symbol table entries for the program in Fig. 2.18 after scanning line 160.

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### One-Pass Assembler Producing Object Code

- Forward reference are entered into the symbol table's list as before
  - If the operand contains an undefined symbol, use 0 as the address and write the Text record to the object program.
- □ However, when definition of a symbol is encountered, the assembler must generate <u>another Text record</u> with the <u>correct operand address</u>.
- □ When the program is loaded, this address will be inserted into the instruction by *loader*.
- □ The object program records must be kept in their original order when they are presented to the loader

## Example

- □ In Fig. 2.20
  - Second Text record contains the object code generated from lines 10 through 40
    - □ The operand addressed for the instruction on line 15, 30, 35 have been generated as 0000
  - When the definition of ENDFIL is encountered
    - □ Generate the third Text record
      - Specify the value 2024 (the address of ENDFIL) is to be loaded at location 201C (the operand field of JEQ in line 30)
      - Thus, the value 2024 will replace the 0000 previously loaded

# Object Program from one-pass assembler for Fig 2.18 (Fig 2.20)

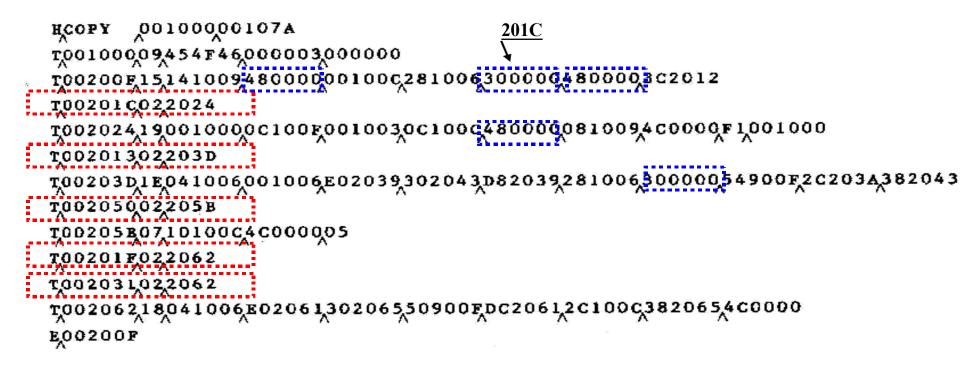
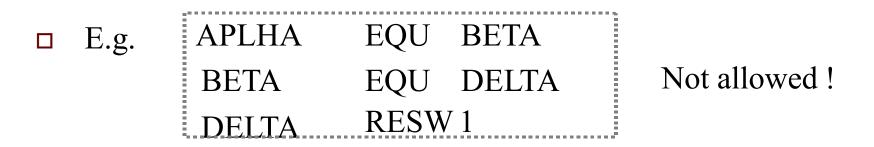


Figure 2.20 Object program from one-pass assembler for program in Fig. 2.18.



## 2.4.2 Multi-Pass Assemblers

- □ Motivation: for a 2-pass assembler, any symbol used on the *right-hand side* should be defined previously.
  - No forward references since symbols' value can't be defined during the first pass



# Multi-Pass Assemblers (Cont.)

- Multi-pass assemblers
  - Eliminate the restriction on EQU and ORG
  - Make as many passes as are needed to process the definitions of symbols.
- □ Implementation
  - To facilitate symbol evaluation, in SYMTAB, each entry must indicates *which symbols are dependent on the values of it*
  - Each entry keeps a <u>linking list</u> to keep track of whose symbols' value depend on an this entry

# Example of Multi-pass Assembler Operation (fig 2.21a)

HALFSZ MAXLEN PREVBT	EQU EQU EQU	MAXLEN/2 BUFEND- BUFFER BUFFER-1
•		
BUFFER	RESB	4096
BUFEND	EQU	*

# Example of Multi-Pass Assembler Operation (Fig 2.21b)

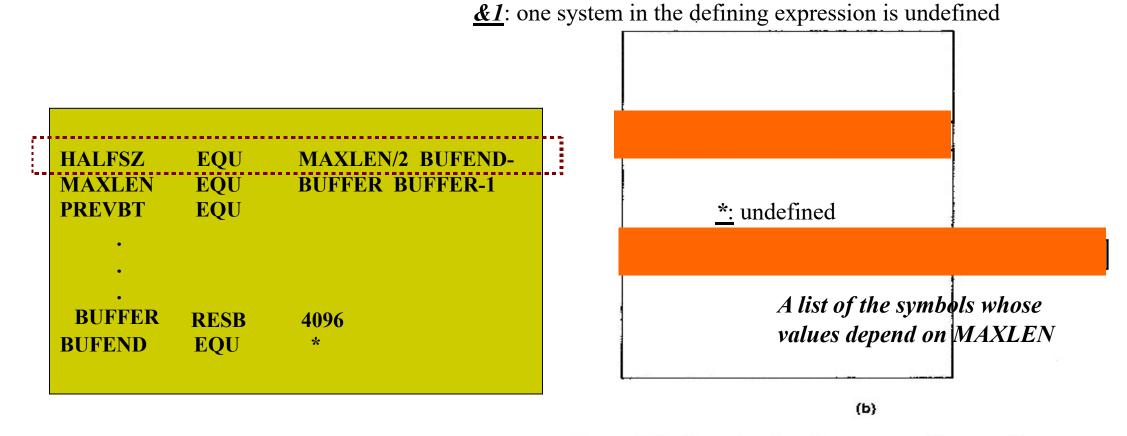
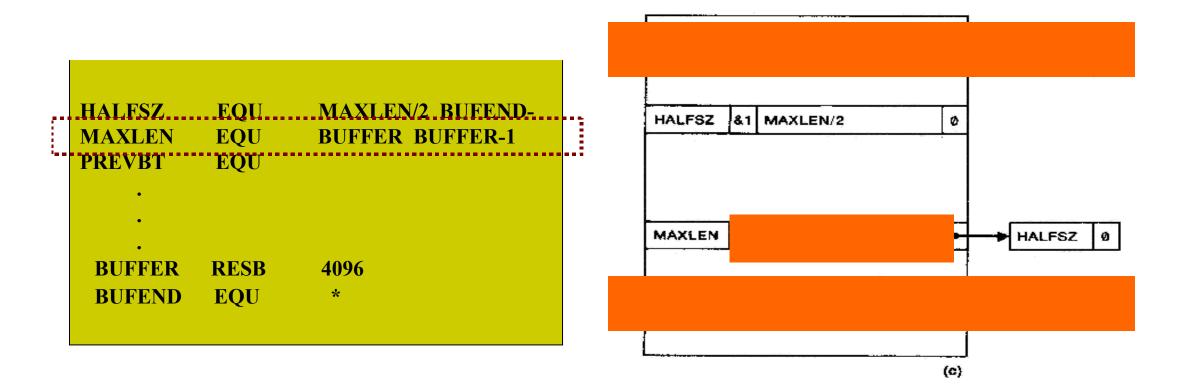
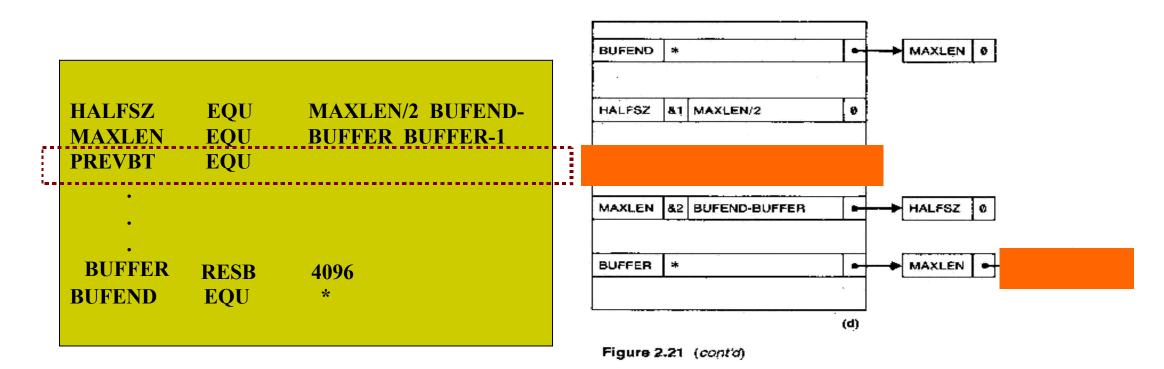


Figure 2.21 Example of multi-pass assembler operation.

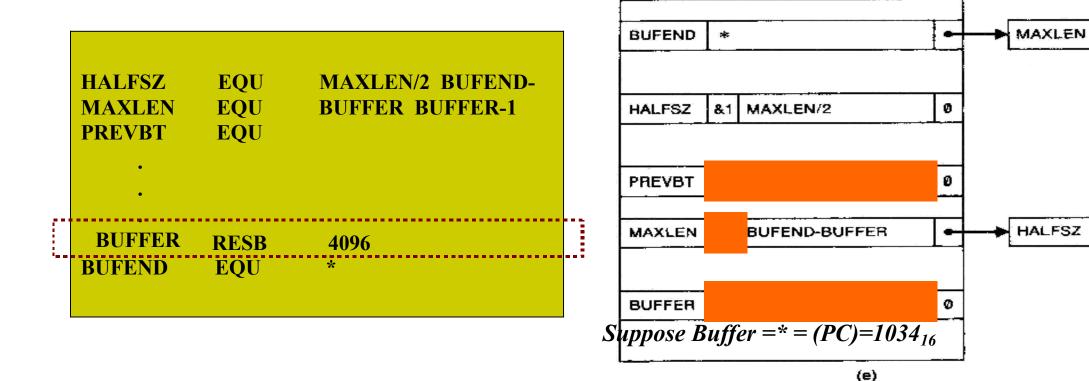
# Example of Multi-Pass Assembler Operation (Fig 2.21c)



# Example of Multi-pass Assembler Operation (fig 2.21d)



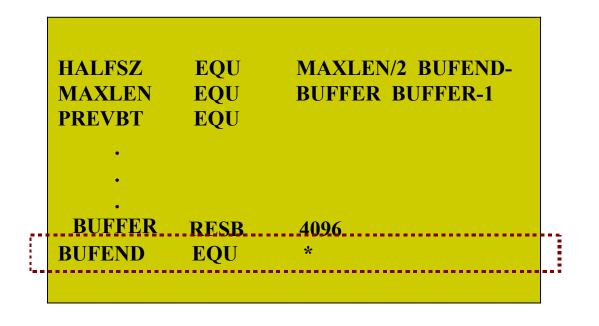
### Example of Multi-pass Assembler Operation (fig 2.21e)



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# Example of Multi-pass Assembler Operation (Fig 2.21f)



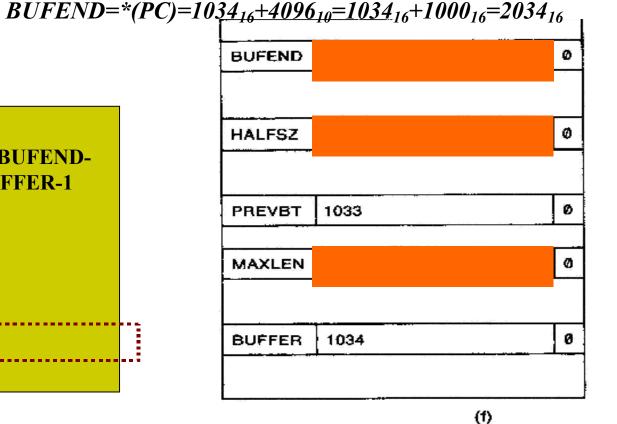


Figure 2.21 (con'd)

# 2.5 Implementation Examples

Microsoft MASM Assembler

Sun Sparc Assembler

□ IBM AIX Assembler

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# 2.5.1 Microsoft MASM Assembler

Microsoft MASM assembler for Pentium and other x86 systems

Programmer of an x86 system views memory as a collection of <u>segments</u>

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If we use a two-pass assembler, the following symbol definition cannot be allowed.

ALPHA EQU BETA

**BETA EQU DELTA** 

DELTA RESW 1

This is because ALPHA and BETA cannot be defined in pass 1. Actually, if we allow multi-pass processing, DELTA is defined in pass 1, BETA is defined in pass 2, and ALPHA is defined in pass 3, and the above definitions can be allowed.

This is the motivation for using a multi-pass assembler.



- It is unnecessary for a multi-pass assembler to make more than two passes over the entire program.
- Instead, only the parts of the program involving forward references need to be processed in multiple passes.
- The method presented here can be used to process any kind of forward references.
- Use a symbol table to store symbols that are not totally defined yet.
- For a undefined symbol, in its entry, We store the names and the number of undefined symbols which contribute to the calculation of its value. – We also keep a list of symbols whose values depend on the defined value of this symbol.
- When a symbol becomes defined, we use its value to reevaluate the values of all of the symbols that are kept in this list.
- The above step is performed recursively.



• Examples

Microsoft MASM Assembler, Sun Sparc Assembler, IBM AIX Assembler

- Microsoft MASM Assembler
- SEGMENT a collection segments, each segment is defined as belonging to a particular class, CODE, DATA, CONST, STACK
- registers: CS (code), SS (stack), DS (data), ES, FS, GS
- similar to program blocks in SIC I ASSUME
- e. g. MOVE ES: DATASEG 2 AX, DATASEG 2 ES, AX » similar to BASE in SIC 11

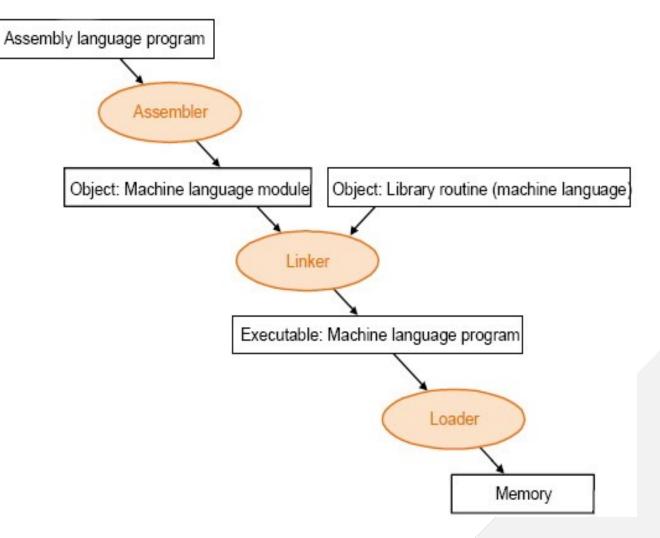


#### • Microsoft MASM Assembler (Contd.)

- JUMP with forward reference
- near jump: 2 or 3 bytes
- far jump: 5 bytes
- e. g. JMP TARGET
- Warning: JMP FAR PTR TARGET
- Warning: JMP SHORT TARGET
- Pass 1: reserves 3 bytes for jump instruction phase error PUBLIC, EXTRN
- similar to EXTDEF, EXTREF in SIC 12



## Advanced Assembly process



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# Advanced Assembly Process (Example)

Mnemonic	Operands	Comment					
MOV	AX,BX	;	Put	byte	count	into	AX

The assembler reads a line like this one from the source code file and writes the equivalent machine instruction to the object code file:



# Advanced Assembly Process

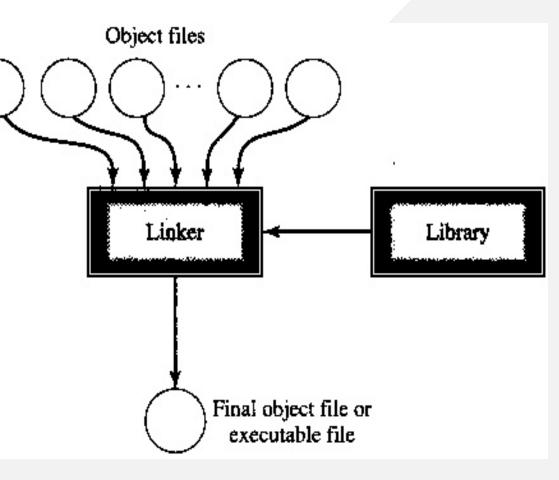
- Assembling
- At assembly time, the assembler:
  - Evaluates conditional-assembly directives, assembling if the conditions are true.
  - Expands macros and macro functions.
  - Evaluates constant expressions such as **MYFLAG AND 80H**, substituting the calculated value for the expression.
  - Encodes instructions and non address operands. For example, **mov cx, 13**; can be encoded at assembly time because the instruction does not access memory.
  - Saves memory offsets as offsets from their segments.
  - Places segments and segment attributes in the object file.
  - Saves placeholders for offsets and segments (relocatable addresses).
  - Outputs a listing if requested.
- Passes messages (such as INCLUDELIB) directly to the linker.



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# Advanced Assembly Process

- Once your source code is assembled, the
  resulting object file is passed to the linker. At
  this point, the linker may combine several
  object files into an executable program. The
  linker:
  - Combines segments according to the instructions in the object files, rearranging the positions of segments that share the same class or group.
  - Fills in placeholders for offsets (relocatable addresses).
  - Writes relocations for segments into the header of .EXE files (but not .COM files).
  - Writes the result as an executable program file.





# Advanced Assembly Process

#### Loading

After loading the executable file into memory, the operating system:

- Creates the program segment prefix (PSP) header in memory.
- Allocates memory for the program, based on the values in the PSP.
- Loads the program.
- Calculates the correct values for absolute addresses from the relocation table.
- Loads the segment registers SS, CS, DS, and ES with values that point to the proper areas of memory.



# Advanced Assembly Process

#### **Useful Tools and Utilities**

- DUMPBIN disassembly program
- Debuggers: OllyDbg and WinDbg
- Consol I/O: iolib.



### References

- [PDF] Systems Programming and Operating Systems by Dhamdhere -Free Download PDF (dlscrib.com)
- [PDF] Principles of Compiler Design By Alfred V. Aho & J.D.Ullman Free Download – Learnengineering.in



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